Decision Diagrams in Combinatorics

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Conference in honor of Tony Guttmann's 80th birthday

Talk Contents

- What is a binary decision diagram (BDD)?
 - How is it useful in combinatorics?
 - How about a other variants?
- Use in pattern avoiding permutations
 - Generalization to multiset MBDD, MZDD, etc.
 - Paper with Tony:

"Counting Occurrences of Patterns in Permutations" January 2025 The Electronic Journal of Combinatorics 32(1) DOI:10.37236/12963

What is a decision diagram?

- A data structure
- Represents a binary function
 - e.g. f(a,b,c) = (a xor b) and (not c)

The Art of Computer Programming, Volume 4, Fascicle 1, Bitwise Tricks & Techniques; Binary Decision Diagrams Donald Knuth

Combinatorics motivation

- A binary function can be considered to define a set S = {x:f(x)=true}
- e.g. variables are bonds on a fixed size lattice, function is true if they form a self avoiding walk.
- Number of true elements = number of self avoiding walks.



Decision Diagram

В

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What can you do with an xDD?

- Often store a complex function of n variables in much less space than O(2ⁿ).
- Logical operations (and or not) of sizes s and t take worst case O(st) and often much better
- Fast to count, find smallest number of satisfying elements, get an arbitrary solution

Application : SAWs

- Self avoiding walk on a finite rectangular lattice
- Each bond is a variable
- Function is true if it is a SAW
- Same algorithm as to count them, just with a more complex answer than an integer.

Application : Tilings

 Knuth suggest the problem of tiling a chessboard with dominoes – how many ways?







Speed

- BDD: 2.1 ms 5679 nodes (11712 to generate)
 - 41 lines of code including comments, not counting xDD library
- ZDD: 3.2 ms 2298 nodes (19790 to generate)
- Simple dynamic programming approach: 6 µs
 - 27 lines of code including comments

Program: https://github.com/AndrewConway/xdd

Bigger problem

- Tiling with unimoes, dominoes, trionimoes:
- 92109458286284989468604
- 243057372832 fewest tiles
- BDD: 1.4s, 2,447,347 nodes
- ZDD: 1.7s, 512,225 nodes



Application : Directed animals

- Polyonimoes on a square lattice.
 - Each occupied site other than the root has to have an element directly below or to the left
 - One variable for each site on a sufficiently large lattice.
 - All constraints are local (suits xDD).



Performance

- 20 terms. 86,574,831 animals
- BDD : 7.8s
 - ⁻ 2,191,868 nodes 23,292,361 used
- ZDD : 6.2s
 - ⁻ 1,292,942 nodes 13,952,387 used
- Dynamic programming :228µs
 - 5022 cache elements

Not looking great for xDDs.

- Simple program much faster for counting.
- xDDs make it easier to make pictures. Possibly more straight forward?

Pattern Avoiding Permutations

- Let p be a permutation of 1..n
- Let P be a permutation of 1...m $m \le n$
- p contains the pattern P if a subset of p of length m has the same relative ordering as P.
 - Otherwise p *avoids* the pattern P.
- Examples:
 - 1,5,3,2,4 contains the pattern 1,3,2 as the subsequence 1,5,3 is in the same relative order as 1,3,2. (also 1,5,2 and 1,5,4 and 1,3,2)
 - 1,2,3,4,5 avoids the pattern 1,3,2 (also avoids 2,1)

Enumeration of pattern avoiding permutations : length 1 to 4

- All but one class solved analytically.
- Dedicated algorithms for enumeration of the 1324 class.

Enumeration of PAPs : length 5

- Too many hard patterns for specialized algorithms
- Best general algorithms explicitly generated all solutions, so slow

Permutation Decision Diagrams

- Minato had a novel idea for representing a set of permutations
 - Represent permutations by a set of binary variables
 - Make a new operation on such ZDDs to compose permutations, maintaining canonicity.
 - $^-$ These are called $\pi DDs.$
- Inoue improved the encoding, called Rot- π DDs.

Use of π DDs in PAPs

- Minato defined a set of operations that created a set of all pattern containing permutations.
- Then subtract from n! to get PAPs.
- Much faster than direct enumeration.
- Inoue's variation is better still.

Pattern containing permutations

- Next step : How many permutations contain a pattern exactly 1 time? Or r times?
- For simple patterns, there are some analytic solutions.
- More generally...?

Minato's algorithm and multisets

- Take Minato's algorithm for generating all pattern containing permutations
- Use multisets instead of sets.

- Where $\{a,b,b,c\} \times \{b,b,c\} = \{b,b,b,b,c\}$

• Now the multiplicity of each element is the number of times it contains the pattern.

Multisets and decision diagrams

- Decision diagrams can be extended to represent multisets.
 - f(booleans) → \mathbb{N}
- Modification: instead of a reference to a node, have an integer multiple and a reference to a node.





Canonicalization

- Need to have a unique representation for each multiset.
 - [–] A multiple of 0 iff the reference is to 0.
 - Each node has two references. Their multiples should be coprime (or 1 if the other is 0).
- Details and theory in section 2 of:

"Counting Occurrences of Patterns in Permutations" January 2025 The Electronic Journal of Combinatorics 32(1) DOI:10.37236/12963

Implementation

- Rust library at https://github.com/AndrewConwa y/xdd
- Can enumerate by multiplicity efficiently.
- Multiset BDD is an MBDD.
 - [–] Similarly MZDD, M π DD,Rot-M π DD
- Program used in paper is in that as an example.

References

• Theory and results

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- Library and program
 - https://github.com/AndrewConway/xdd
- Program to make pictures & benchmarks in this talk
 - https://github.com/AndrewConway/xdd_pictures