A domain wall theory for the prioritising exclusion process

Caley Finn in collaboration with Jan de Gier

University of Melbourne

December 2012

Customers waiting to be served:

(1) (2) (1) (2) (1) (1)

▶ Priority class high (1) or low (2)

Customers waiting to be served:



- ▶ Priority class high (1) or low (2)
- Arrive at rate λ_i customers / minute

Customers waiting to be served:

$$\begin{array}{c} \overbrace{1}^{\lambda_{1}} \\ \overbrace{2}^{\lambda_{2}} \end{array} \begin{array}{c} 1 \end{array} \begin{array}{c} 2 \end{array} \begin{array}{c} 1 \end{array} \begin{array}{c} 2 \end{array} \begin{array}{c} 1 \end{array} \begin{array}{c} \mu \end{array} \end{array} \begin{array}{c} \mu \end{array} \end{array} \begin{array}{c} 1 \end{array}$$

- ▶ Priority class high (1) or low (2)
- Arrive at rate λ_i customers / minute
- Served at rate μ customers / minute

Customers waiting to be served:



- ▶ Priority class high (1) or low (2)
- Arrive at rate λ_i customers / minute
- Served at rate μ customers / minute
- ► High priority overtake low at rate *p* places / minute

Motivation

A related model: the *accumulating priority queue*. Customers ordered according to priority

$$V = b_i(t - t_{\mathsf{arrive}}), \quad i = 1, 2, \quad b_1 \ge b_2$$



The Prioritising Exclusion Process

The queue is equivalent to an exclusion process - the PEP



- Low priority \rightarrow empty lattice site,
- High priority \rightarrow filled lattice site
- $\blacktriangleright \text{ Overtaking} \rightarrow \text{particle hopping}$
- But, lattice length not fixed
- Specify a configuration by

$$\boldsymbol{\tau} = (\tau_n, \tau_{n-1}, \ldots, \tau_1)_n, \qquad \tau_i \in \{0, 1\}$$

Bounded and unbounded queues

Let
$$\lambda = \lambda_1 + \lambda_2$$
.
 $\blacktriangleright \lambda > \mu$
 $\langle n \rangle \sim (\lambda - \mu)t$ expected length is unbounded
 $\flat \lambda < \mu$
 $\langle n \rangle = \frac{\lambda}{\mu - \lambda}$ expected length is bounded
 $P_n = \left(1 - \frac{\lambda}{\mu}\right) \left(\frac{\lambda}{\mu}\right)^n$

 P_n is the stationary length distribution, solution to

$$0 = \frac{d}{dt}P_0 = \mu P_1 - \lambda P_0$$

$$0 = \frac{d}{dt}P_n = \lambda P_{n-1} + \mu P_{n+1} - (\lambda + \mu)P_n, \quad n > 0$$

Density profiles

• Density by position (i) and queue length (n)

 $\langle \tau_i \rangle_n = P(HI \text{ in place } i \text{ and queue length is } n)$

Simulation results show two distinct phases



The domain wall model is an idea borrowed from the ASEP



 Domain wall between *jam* and low density region located at first empty site from service end

The domain wall model is an idea borrowed from the ASEP



- Domain wall between *jam* and low density region located at first empty site from service end
- ► Jam grows when particle jumps on to the end

The domain wall model is an idea borrowed from the ASEP



- Domain wall between *jam* and low density region located at first empty site from service end
- ► Jam grows when particle jumps on to the end
- Jam shrinks when service occurs

Assume the region beyond the jam has constant low density lpha

$$P(\tau_{i} = 1 | n, k)$$

$$(\tau_{i} = 1 | n, k) = \begin{cases} 1, & 1 \le i \le k \\ 0, & i = k + 1 \\ \alpha, & k + 2 \le i \le n \end{cases}$$

P(n, k), probability of queue length n, jam length k, should satisfy

$$\frac{d}{dt}P(n,k) = \lambda P(n-1,k) + p\alpha P(n,k-1) + \mu P(n+1,k+1) - (\lambda + \mu + p\alpha)P(n,k)$$

For n > k + 1, n, k > 0.

Assume the region beyond the jam has constant low density lpha

$$\begin{array}{c|c} \lambda & & & & & \\ \hline & & & & & \\ \hline & & & & \\ n & & & & \\ \hline n & & \\ n & & \\ \hline n & & \\ n & & \\ \hline n & & \\ n & & \\ n & & \\ \hline n & & \\ n &$$

P(n, k), probability of queue length n, jam length k, should satisfy

$$\frac{d}{dt}P(n,k) = \lambda P(n-1,k) + p\alpha P(n,k-1) + \mu P(n+1,k+1) - (\lambda + \mu + p\alpha)P(n,k)$$

For n > k + 1, n, k > 0.

Simple domain wall model - boundary conditions

- Simple domain wall model is complicated by the boundary conditions – what happens when the jam reaches the arrival end?
- Simplification: consider the $n \to \infty$ limit

$$\frac{d}{dt}P(k) = p\alpha P(k-1) + \mu P(k+1) - (\mu + p\alpha)P(k)$$

• The only boundary case is k = 0

The unbounded queue – an exact solution

Stationary probability of finite segment from an infinite queue:

$$P(\tau_m, \tau_{m-1}, \ldots, \tau_1) = \sum_{\tau_{\infty}, \ldots, \tau_{m+1} = 0, 1} P(\ldots, \tau_{m+1}, \tau_m, \tau_{m-1}, \ldots, \tau_1)$$

Example with length k jam

Exact equation:

$$0 = \frac{d}{dt} P(1, 0, 1, 0, 1^{k})$$

= $pP(\mathbf{1}, \mathbf{0}, 0, 1, 0, 1^{k}) + pP(1, \mathbf{1}, \mathbf{0}, 0, 1^{k}) + pP(1, 0, 1, \mathbf{1}, \mathbf{0}, \mathbf{1^{k-1}})$
+ $\mu P(1, 0, 1, 0, \mathbf{1^{k+1}}) + \mu P(1, 0, 1, 0, 1^{k}, \mathbf{0})$
- $(\mu + 2p)P(1, 0, 1, 0, 1^{k})$

Domain wall ansatz



Applied to the example

$$0 = \frac{d}{dt} P_{jam}(k)$$

= $p \alpha P_{jam}(k-1) + \mu P_{jam}(k+1) + \mu (1-\alpha) \alpha^k P_{jam}(0)$
 $- (\mu + p \alpha) P_{jam}(k)$

Almost identical to the simple domain wall dynamics

Recurrence for $P_{jam}(k)$

Combining with k = 0 case gives the recurrence

$$P_{\mathrm{jam}}(k) = rac{plpha}{\mu} P_{\mathrm{jam}}(k-1) + lpha^k P_{\mathrm{jam}}(0), \qquad k \geq 1$$

Solution

$$\begin{split} P_{\mathrm{jam}}(k) &= \sum_{i=0}^{k} \left(\frac{p\alpha}{\mu}\right)^{k-i} \alpha^{i} P_{\mathrm{jam}}(0) \\ &= (p\alpha)^{k} \frac{\left(\frac{1}{p}\right)^{k+1} - \left(\frac{1}{\mu}\right)^{k+1}}{\frac{1}{p} - \frac{1}{\mu}} P_{\mathrm{jam}}(0) \end{split}$$

 $P_{\rm jam}(0)$ determined by the normalisation condition

$$\sum_{k=0}^{\infty} P_{\rm jam}(k) = 1$$

Constraint $p\alpha < \mu$: domain wall stays near service end

The arrival frame

Reference frame at the arrival end of the queue

For a finite segment at the arrival end, rate equation includes arrival terms

$$0 = \frac{d}{dt} P(\tau_1, \tau_2, \dots, \tau_m)$$

= $\dots \tau_1 \lambda_1 P(\tau_2, \dots, \tau_m) + (1 - \tau_1) \lambda_2 P(\tau_2, \dots, \tau_m)$
 $- \lambda P(\tau_1, \dots, \tau_m) + \dots$

The arrival frame solution

Assuming the jam stays far from the arrival end, the *arrival frame* ansatz is

$$P(\tau_1,\ldots,\tau_m)=\alpha^{m_1}(1-\alpha)^{m-m_1}$$

All cases reduce to

$$p\alpha^{2} - (p + \lambda)\alpha + \lambda_{1} = 0$$

$$\Rightarrow \alpha = \frac{p + \lambda - \sqrt{(p - \lambda)^{2} + 4p\lambda_{2}}}{2p}$$

Note $p\alpha < \lambda$: domain wall stays far from arrival end

Exact solution vs Monte Carlo simulation

Service frame density from domain wall solution

$$egin{aligned} &\langle au_i
angle &= lpha \sum_{k=0}^{i-2} P_{\mathrm{jam}}(k) + 0 imes P_{\mathrm{jam}}(i-1) + 1 imes \sum_{k=i}^{\infty} P_{\mathrm{jam}}(k) \ &= lpha + (1-lpha) \left(rac{p lpha}{\mu}
ight)^i \end{aligned}$$

• Monte Carlo simulation, $\lambda_1 = 1.3$, $\lambda_2 = 0.2$, $\mu = 1$, $p = 0.3, 1, 1.5 \rightarrow p\alpha \simeq .25, .74, .95$



The bounded queue

Density profiles have strong dependence on the length distribution

►
$$\langle \tau_i \rangle_n \leq P_n, \qquad P_n = \left(1 - \frac{\lambda}{\mu}\right) \left(\frac{\lambda}{\mu}\right)^n$$

Can we factor out length dependence?



Domain wall ansatz with length assumption

$$P(\tau_n, \tau_{n-1}, \dots, \tau_{k+2}, 0, 1^k; n) = P_n \alpha^{n_1} (1-\alpha)^{n-k-n_1-1} P_{jam}^*(k)$$

Length assumption solution

Solution for domain wall away from boundaries

$$egin{aligned} \mathcal{P}^*_{ ext{jam}}(k) &= \sum_{i=0}^k \left(rac{plpha}{\lambda}
ight)^{k-i} lpha^i \mathcal{P}^*_{ ext{jam}}(0) \ &= (plpha)^k rac{\left(rac{1}{p}
ight)^{k+1} - \left(rac{1}{\lambda}
ight)^{k+1}}{rac{1}{p} - rac{1}{\lambda}} \mathcal{P}^*_{ ext{jam}}(0), \end{aligned}$$

for $p\alpha < \lambda$, and

$$lpha = rac{m{p} + \mu - \sqrt{(m{p} + \mu)^2 - 4m{p}rac{\lambda_1\mu}{\lambda}}}{2m{p}}$$

Solution fails if the jam reaches the arrival end, e.g.

$$\frac{d}{dt}P(0,1^{n-1};n)\neq 0$$

An approximate solution to an approximation of the problem \rightarrow compare to simulations

Approximate solution vs Monte Carlo simulation

With
$$\mu = 1, p = 1, \lambda_1 = 0.7, \lambda_2 = 0.2 \Rightarrow p\alpha = 0.53$$



When $p\alpha < \lambda$, domain wall solution with length assumption (solid lines) gives a very good approximation.

Monte Carlo results for $p\alpha > \lambda$

With $\mu = 1, p = 3, \lambda_1 = 0.7, \lambda_2 = 0.1 \Rightarrow p\alpha = 0.83$



For $p\alpha > \lambda$, the domain wall solution does not apply and even the length assumption is clearly invalid.

Stationary phase diagram

- Fixed λ , and $\mu = 1$
- Phase diagram for $0 \le \lambda_1 \le \lambda$, p > 0

Unbounded queue ($\lambda = 1.5$)

Bounded queue ($\lambda = 0.6$)





$$\lambda_1 < \min\{1 + \frac{\lambda - 1}{p}, \lambda\}$$



 $p\alpha < \lambda \Rightarrow$

 $\lambda_1 < \min\{\lambda^2\left(1+\frac{1-\lambda}{p}
ight),\lambda\}$

Conclusion

From a simple physical picture, the domain wall idea, we

- ► Found the exact solution for the unbounded queue;
- ► Found a good approximation for the bounded queue;
- Characterised the stationary behaviour, in both cases, by the rate at which the jam of high priority customers grows.

Future work:

- Compute waiting times for the PEP and compare to the accumulating priority queue (APQ).
- Is the change in stationary behaviour we see in the PEP reflected in the APQ?
- ► Can we find an exact solution for the bounded PEP?

The PEP is related to the *accumulating priority queue*: customers ordered according to priority

$$V = b_i(t - t_{arrive})$$



The PEP is related to the *accumulating priority queue*: customers ordered according to priority

$$V = b_i(t - t_{arrive})$$



The PEP is related to the *accumulating priority queue*: customers ordered according to priority

$$V = b_i(t - t_{arrive})$$



The PEP is related to the *accumulating priority queue*: customers ordered according to priority

$$V = b_i(t - t_{arrive})$$

